RAEAX

Rowhammer Amplification by Execution of Additional X86 instructions

Martin Heckel June 29, 2021

Hof University, University of Applied Sciences

Outline

Introduction

Background

Toolset Hammertinger

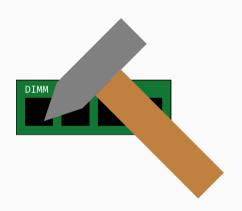
Amplification of rowhammer

Conclusion

Introduction

What is Rowhammer?

- Flipping Bits in Memory Without Accessing Them [3]
- Current DRAM has a very high integration density
- Frequent access to memory locations leads to bit flips in nearby memory locations
- Not the intended behaviour (this should not happen)



What is Rowhammer?

· Flipping Bits in Memory Without Accessing Them So what? Currel **11** Space [...] is **big**. Really big. You densi. just won't believe how vastly hugely • Frequ mindbogglingly big it is. " to bit — Douglas Adams · Not th happen)

Mostly harmless

Local Denial of Service (DoS)

Mostly harmless ... or ...

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- ..

- Local Denial of Service (DoS)
- Esca Problem detected!
- Esca
- · Usir
- · Ove

Let's mitigate it!

- Brea
- ...

Is this still a problem?

- DDR2 systems are typically not affected due to lower integration density
- DDR3 and DDR4 systems got BIOS updates with mitigations starting in 2015
- DDR4 memory modules often contain mitigations additionally.
- · Rowhammer is not a problem anymore

Is this still a problem?

· DDR2 systems are typically not affected due to lower integration density

. DDF Of course, the problem is not solved!

2015

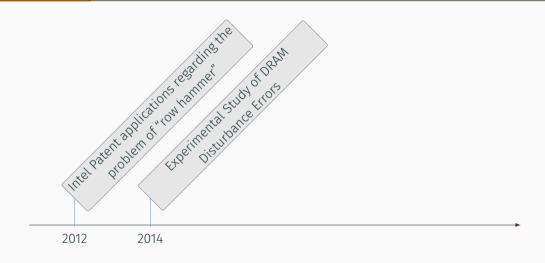
· DDF

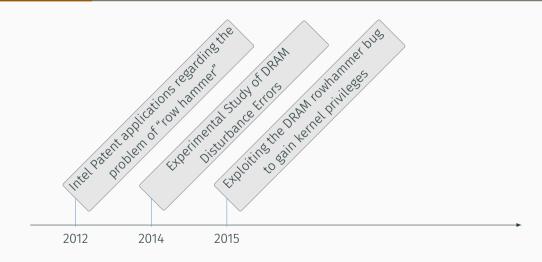
STILL FLIPPING ANYWAY

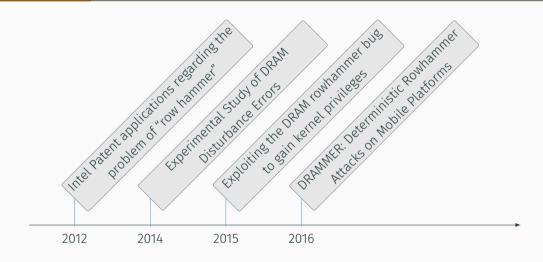
Is this still a problem?

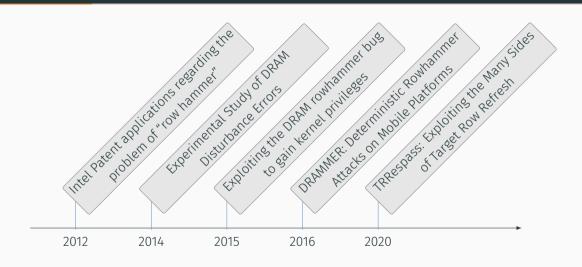
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- DDR3 and DDR4 systems got BIOS updates with mitigations starting in 2015
- DDR4 memory modules often contain mitigations additionally.
- The mitigations on DDR3 systems are not as effective as assumed
- $\frac{1}{4}$ of DDR4 modules are still vulnerable [1]











Background

Reminder: Virtual Memory Management

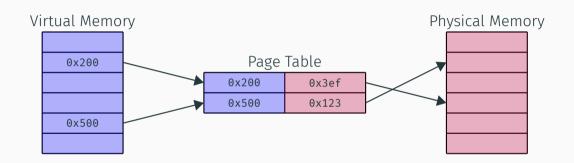
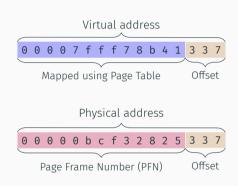


Image from the slides "0x0D Low-Level Fundamentals"

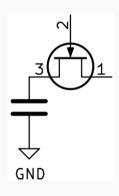
Virtual Memory Management Part 1 — Pages

- Memory is managed in Pages
- On x86 Linux, one page has a size of 4 KiB (12 bit required for addressing)
- An address in virtual memory can be considered as a concatenation of "page address" and "offset".
- First 52 bits are mapped, last 12 bits are "copied"



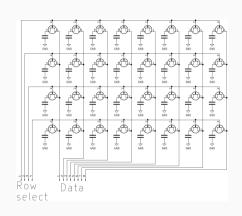
DRAM basics Part 1 — Cells

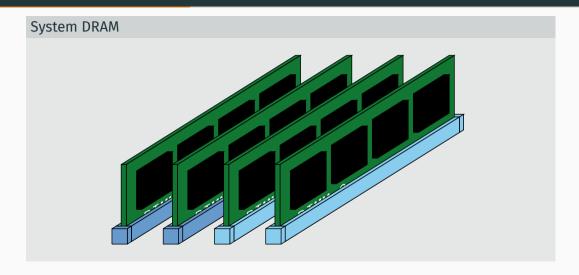
- A single cell consists of a capacitor storing the actual data and a transistor controlling the access
- In this example: control at pin 2 and access at pin 1
- Reading procedure: Enable the control pin and read the voltage at the access pin
- Writing procedure: Apply the level that should be written to the access pin and enable the control pin

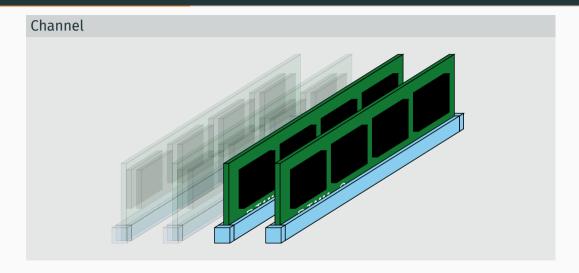


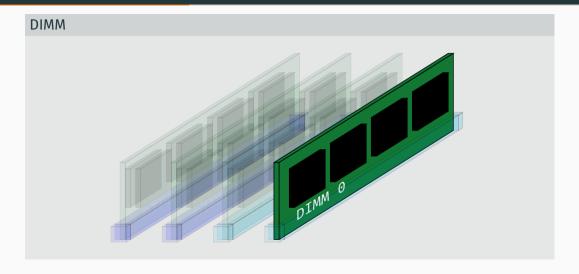
DRAM basics Part 2 — Array

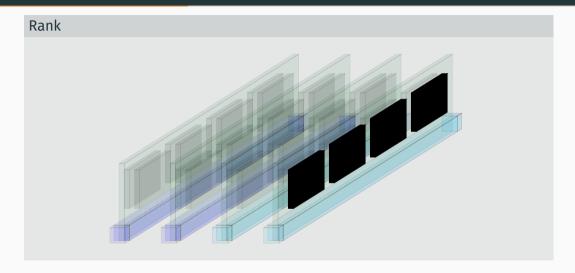
- Multiple of these cells are organized in an array
- Control pins of the cells are connected in rows (only entire rows can be enabled)
- Access pins of the cells are conneted in columns
- Capacitors loose chage over time, so it is required to refresh the cells periodically

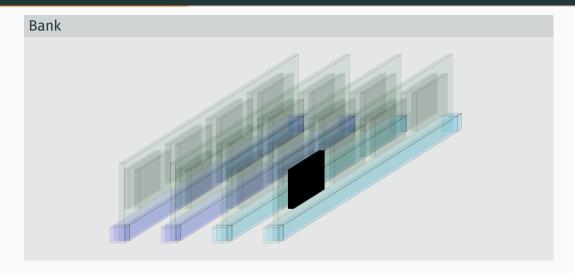


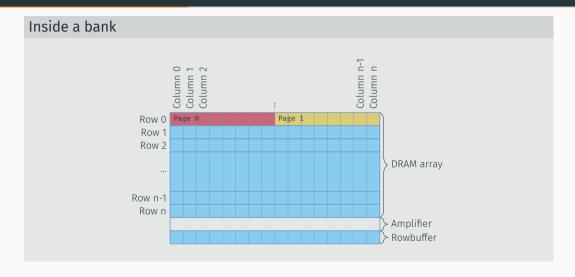












Virtual Memory Management Part 2 — DRAM addressing

- Data is stored in physical memory:
 - · Channel
 - DIMM
 - Rank
 - · Bank
 - Row
 - Column
- The Memory Management Unit (MMU) translates physical addresses to memory locations



Virtual Memory Management Part 2 — DRAM addressing

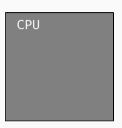
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Rowhammer from the perspective of the CPU

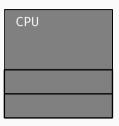
```
hammer:
mov eax, X
mov ebx, Y
clflush X
clflush Y
jmp hammer
```





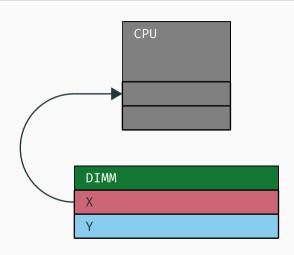
Source code from Wikipedia

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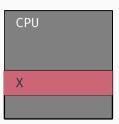




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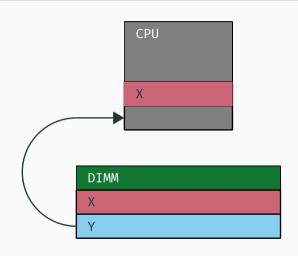


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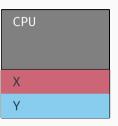




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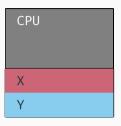


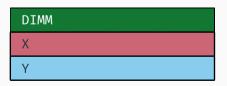
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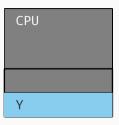


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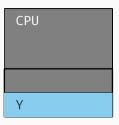


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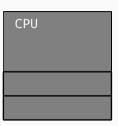


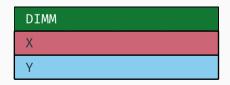
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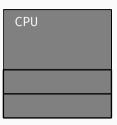


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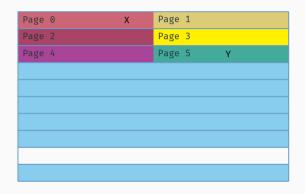


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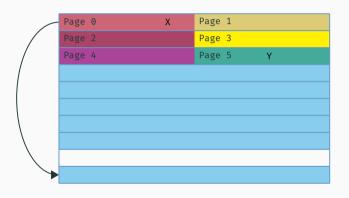
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Page 0	Х	Page 1		
Page 2		Page 3		
Page 4		Page 5	Υ	

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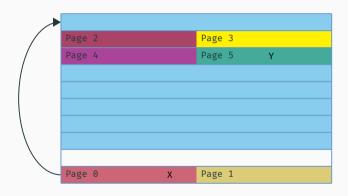
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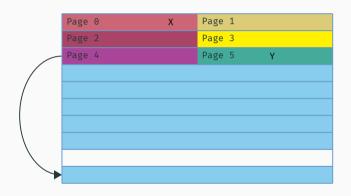
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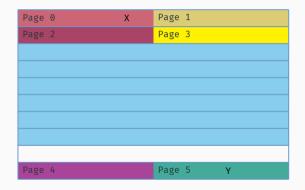
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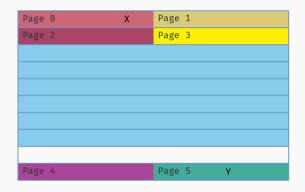
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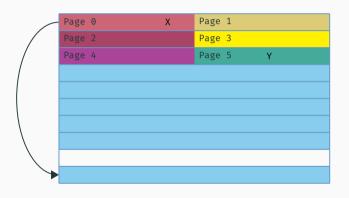
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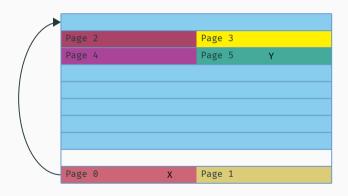
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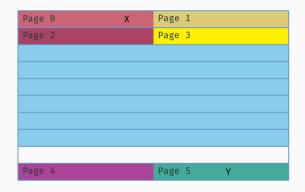
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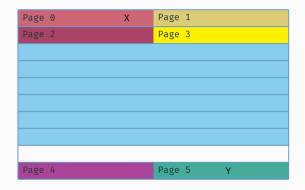
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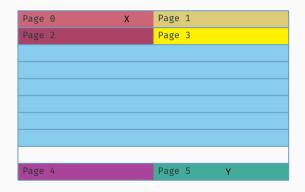
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Rowhammer from the perspective of a DRAM bank (abstract)

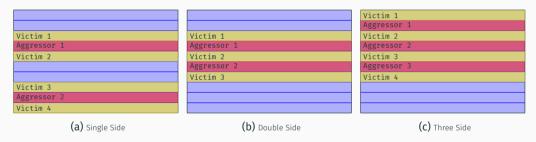
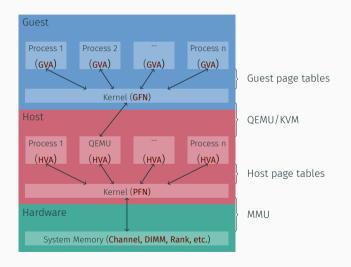


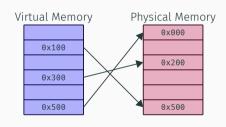
Figure 1: Examples of rowhammer patterns

Virtual Memory Management Part 3 — Virtual environments



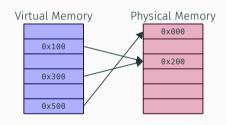
Kernel Samepage Merging (KSM)

- Mechanism in the Linux Kernel
- Deduplicates pages that are marked as merge candidates:
 - Multiple virtual pages with same content are stored once in physical memory
 - Page tables of all pages are adjusted so they reference to the same physical page
 - Copy On Write (COW) policy: Writing to a page copies the page, copy is accessed



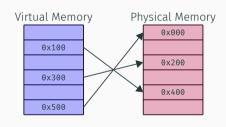
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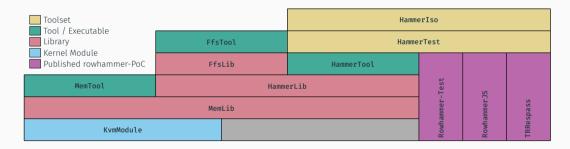
Flip Feng Shui (FFS)

- KSM merges pages across processes
- By default, QEMU/KVM marks all pages as KSM merge candidates
- · When KSM is running, pages across VMs are deduplicated
- Idea from Razavi et al. [4]:
 - Select a target page
 - Chose a bit that should flip inside the target page
 - Find a location in memory that is vulnerable to rowhammer where the specified bit flips
 - Get this target page merged at a location found before
 - · Execute the rowhammer attack at that location
 - The bit should be flipped at all pages that were merged to that location

Toolset Hammertinger

What is Hammertinger?

· All-in-One rowhammer and FFS exploitation framework



Practical rowhammer exploitation

- · At first, no bit flips were found on the tested systems
- · Reason: The systems had a BIOS version with a mitigation for rowhammer
- · Idea: Downgrade BIOS to a version before the mitigation
- Problem: Downgrade locks (unable to perform the downgrade)

Practical rowhammer exploitation

Solution

· At f

· Rea

• Idea

· Pro



Manually downgrade the BIOS

nmer

Demo 1

Exploiting rowhammer with HammerTool

Evaluation of HammerTool

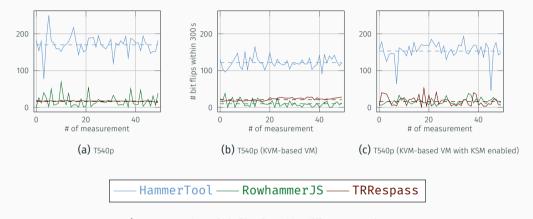


Figure 2: Number of bit flips found by different PoCs in 300s

Amplification of rowhammer

Increase the amount of bit flips with KSM

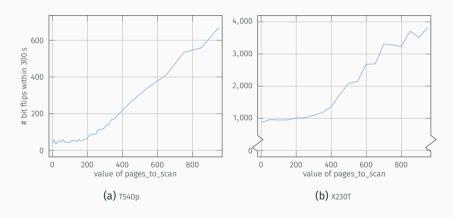


Figure 3: Number of bit flips found by HammerTool in dependence of pages_to_scan

Increase the amount of bit flips with KSM

Increase the amount of bit flips with Flipper

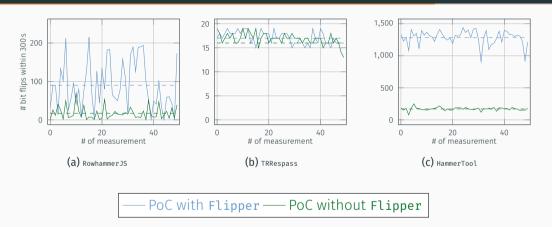


Figure 4: Number of bit flips found by different PoC in 300s with and without Flipper

Demo 2

Increase the amount of bit flips with Flipper

Increase of bit flips with Flipper on a mitigated system

Reference	DIMM	without Flipper	with Flipper	ratio
MO	Samsung M471B5273DH0-CH9	$3.01 \cdot 10^{1}$	$1.8724 \cdot 10^3$	62.2
M1	Samsung M471B5273DH0-CH9	$1.384 \cdot 10^{1}$	$8.157 \cdot 10^{2}$	58.93
M2	Samsung M471B5273DH0-CH9	4.18 · 10 ⁰	$2.5348 \cdot 10^{2}$	60.64
МЗ	Samsung M471B5273DH0-CH9	$9.2 \cdot 10^{-1}$	$3.61 \cdot 10^{1}$	39.23
M4	Samsung M471B5273DH0-CH9	6.54 · 10 ⁰	$3.7384 \cdot 10^{2}$	57.16
M5	Kingston KVR16S11/4 99U5428-049.A00LF	$0 \cdot 10^{0}$	$0 \cdot 10^{0}$	
M6	Samsung M471B1G73BH0-YK0	$2 \cdot 10^{-2}$	$1.52 \cdot 10^{0}$	76

Table 1: Number of bit flips found by HammerTool with and without Flipper

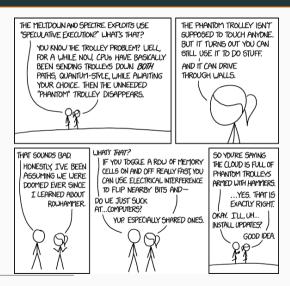


Conclusion

Conclusion

- · An attacker can flip bits in DRAM by accessing other memory locations
- There are sophisticated exploits using rowhammer
- · Mitigation strategies do often not prevent rowhammer but make it unlikely
- HammerTool and Flipper increase the amount of bit flips found on a system significantly which can help bypassing the mitigation on DDR3 systems (maybe, DDR4 as well)

Conclusion



Questions?

References i

- [1] Pietro Frigo et al. "TRRespass: Exploiting the Many Sides of Target Row Refresh". In: S&P. Best Paper Award. May 2020. URL: https://download.vusec.net/papers/trrespass_sp20.pdf.
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